Cloaca

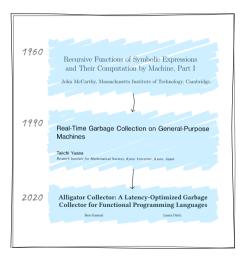
A Concurrent Hardware Garbage Collector for Non-strict Functional Languages

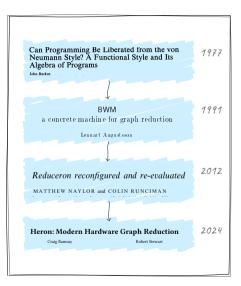
Craig Ramsay & Rob Stewart

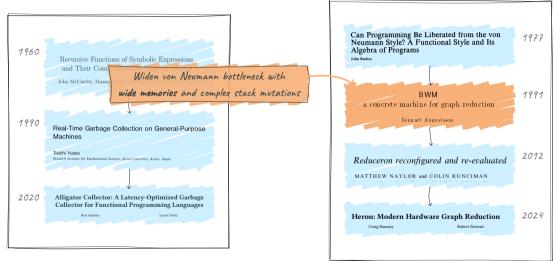
September 2024

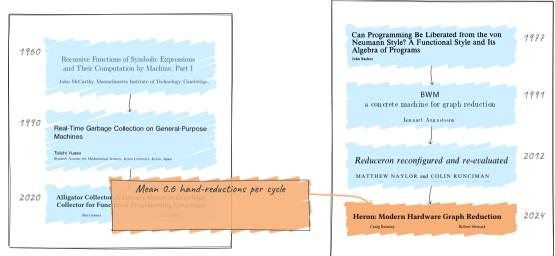
Heriot-Watt University

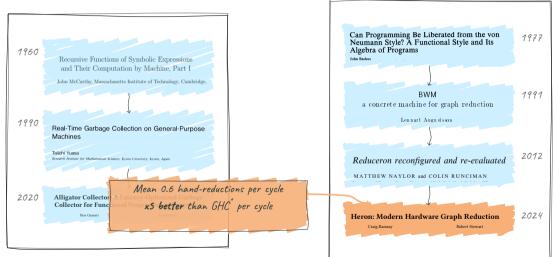


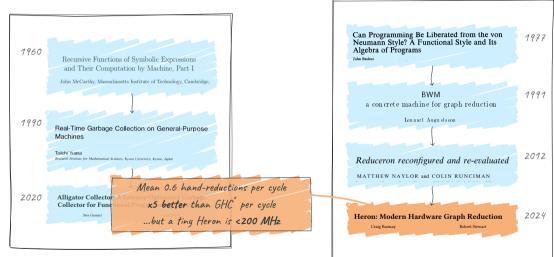


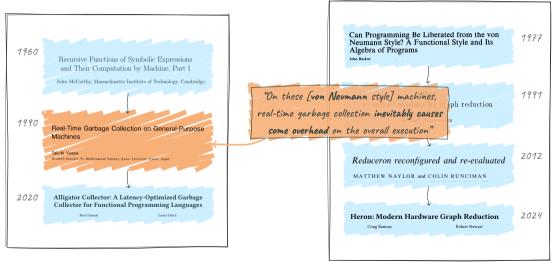


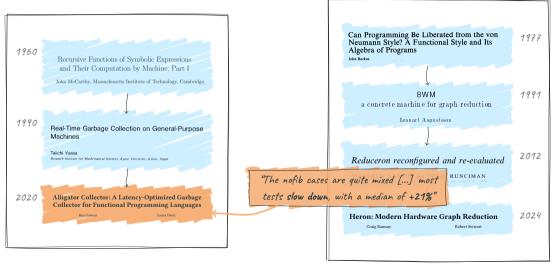


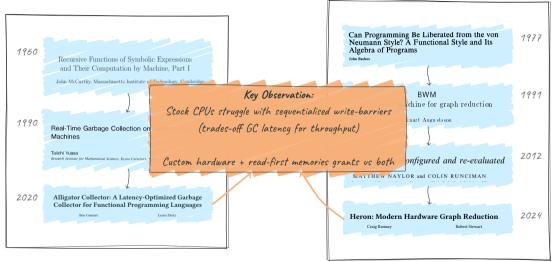








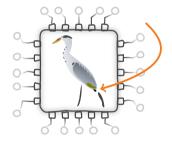




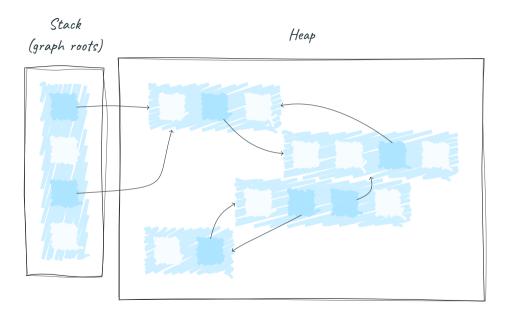
Cloaca

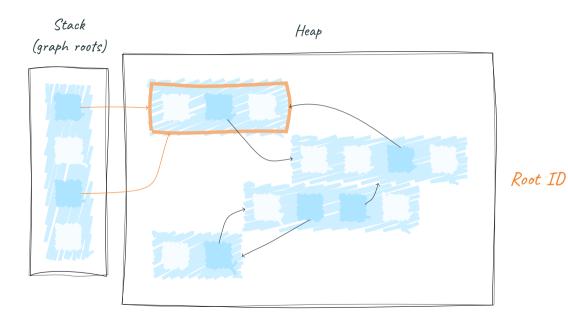
noun [C] /kloh-ah-kuh/

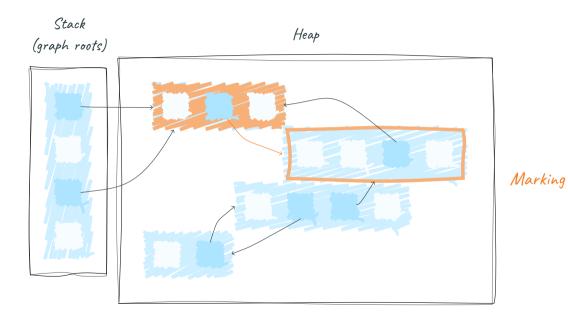
The system responsible for all the waste generated by a heron

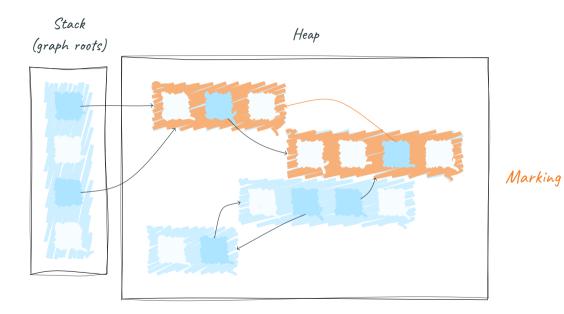


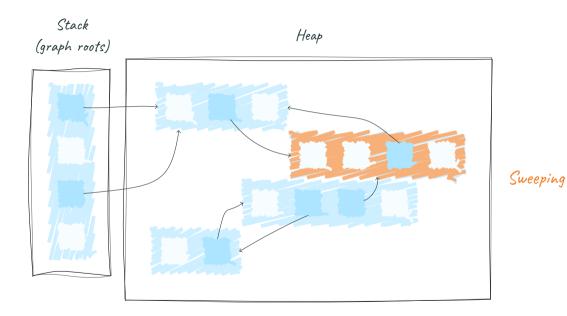


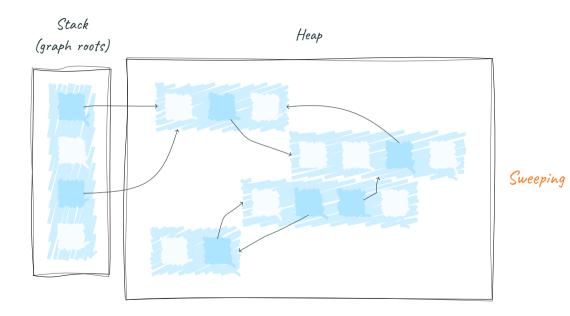


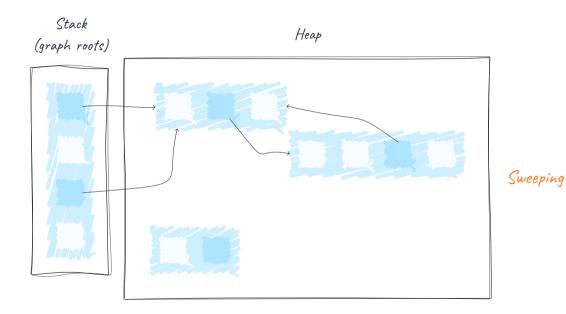


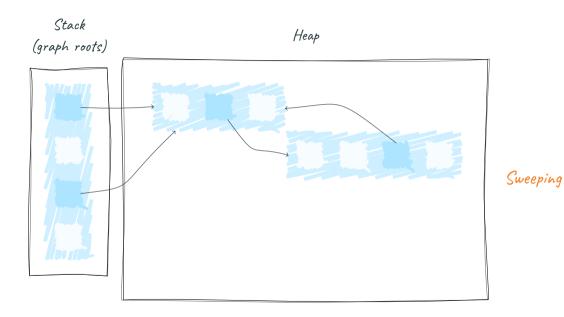


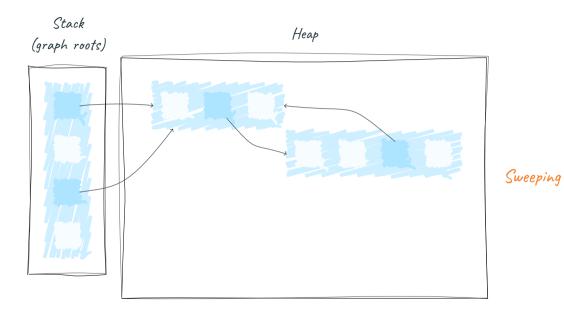












Challenges for concurrent software implementation?

Expect 21% median slowdown for nofib!

1) Allocation depends on GC state

Stop-the-world GC

vs

Concurrent GC

Function alloc (app): if allocBarrier(gcPhase, hp) then tag hp as Marked else tag hp as Unmarked heap[hp] — app hp++

2) Non-moving GC needs complex allocation

Stop-the-world GC

VS

Concurrent GC

Function alloc (app): a ← pop from freelist if allocBarrier(gcPhase, a) then L tag a as Marked else L tag a as Unmarked heap[a] ← app

3) Prevent graph updates from destroying edges

Stop-the-world GC

VS

Concurrent GC

Additional hardware-enabled optimisations

Heron's existing dynamic update avoidance system...

data Atom = ... | Var Shared Int | Arg Shared Int

...

Function unwind (a, shared):

•••

•••

•••

if shared and not NF then push a onto update stack

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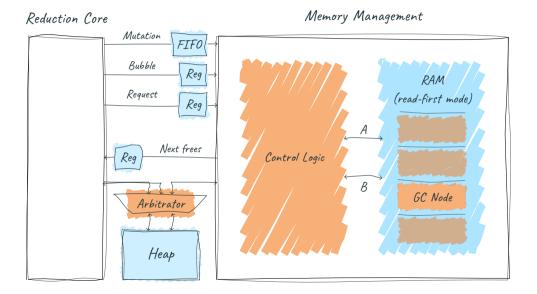
•••

if shared and not NF *then*_ push a onto update stack

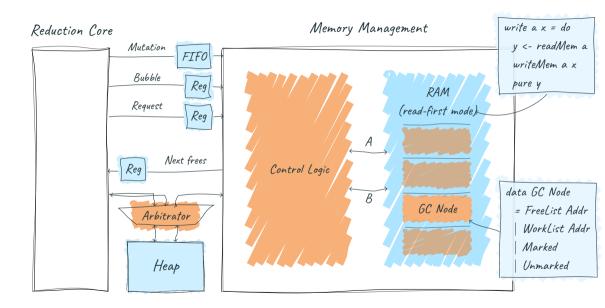
if not shared then ∟ dealloc a

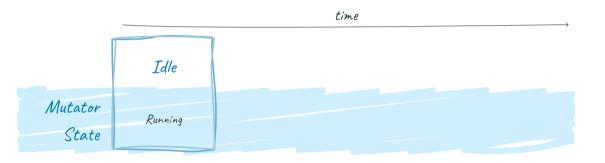
... is just one-bit reference counting with a hat on.

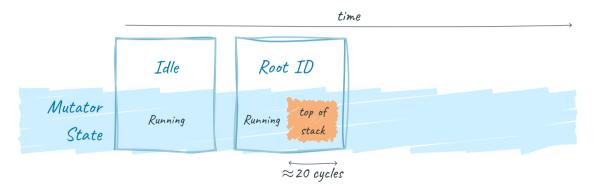
Architecture

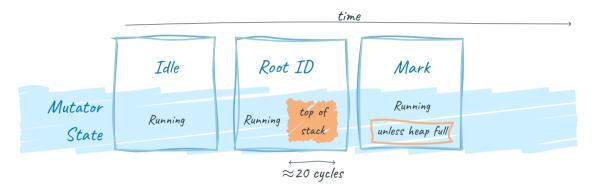


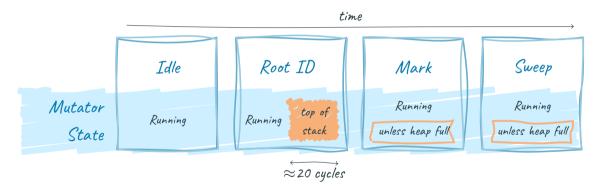
Memory Management Reduction Core Mutation FIFO Bubble Reg RAM Request (read-first mode) Reg А Next frees Control Logic Reg data GC Node В GC Node = FreeList Addr Arbitrator WorkList Addr Marked Неар Unmarked











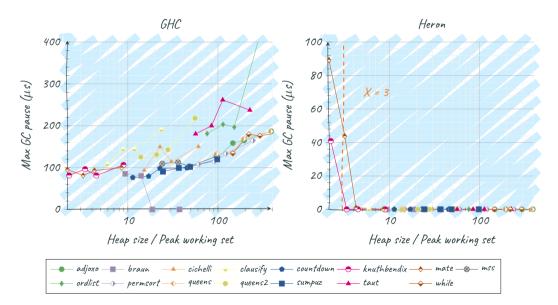
Results

	Peak GHC working set (KB)	GHC Allocations (MB)	LoC
Adjoxo	47	301	72
Braun	46	0	26
Cichelli	52	41	123
Clausify	77	364	69
Countdown	46	54	62
Knuthbendix	105	54	324
Mate	137	430	293
Mss	46	359	13
Ordlist	46	984	28
Permsort	46	2320	10
Queens	55	1038	25
Queens2	47	1084	20
Sumpuz	48	1293	72
Taut	47	236	37
While	46	264	89

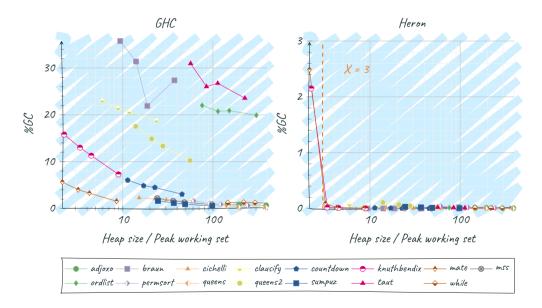
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GC worst-case pause



GC wall-clock overhead



Concurrent tracing $GC \rightarrow$ high throughput and low latency?

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Custom hardware with dual-port read-first memories can handle them in a single cycle.

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Stock processors + software struggle to maintain throughput. Write-barriers and friends are hard.

Custom hardware with dual-port read-first memories can handle them in a single cycle.

Cloaca often pauses for only \approx 20 cycles for a GC pass, and catches \approx 50% of all our garbage before tracing.

Questions?

Break glass in case of emergency

tails [] = [] tails (x : xs) = (x : xs) : tails xs

inits xs =
case xs of
[]
$$\rightarrow$$
 [[]]
(y : ys) \rightarrow xs : inits (init xs)

segments xs = concatMap tails (inits xs)

 $mss = maximum \cdot map sum \cdot segments$

main = let x = 0 - 50 y = 150 in mss \$ enumFromTo x y

e ::==	Expressions	
ē	(Application)	
case e of ā	(Case expression)	
let b in e	(Let expression)	
n	(Integer)	
×	(Variable)	
\otimes	(Primitive Op)	
F	(Function)	
<i>K</i>	(Constructor)	
a ::= K $\bar{x} ightarrow e$	Case alternative	
$b ::= x \mapsto e$	Let binding	
$d ::= f \bar{x} = e$	Function definition	

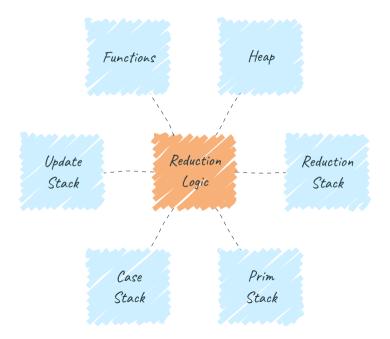
Heron

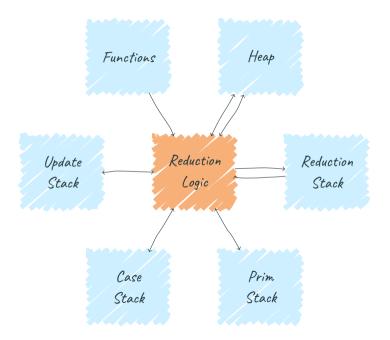
noun [C] /"heran/

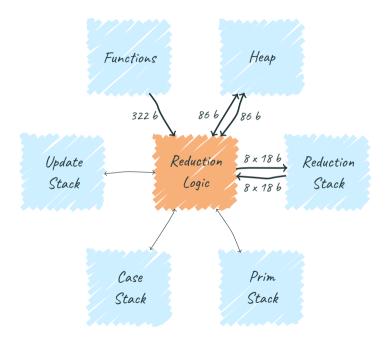
A processor for lazy functional languages. Performs beta reduction in one clock cycle via multiple, wide, multi-ported memories.

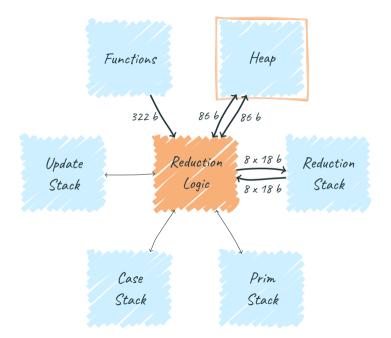


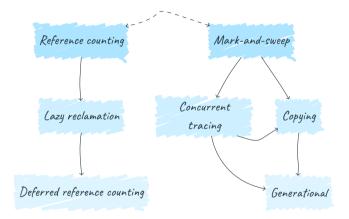


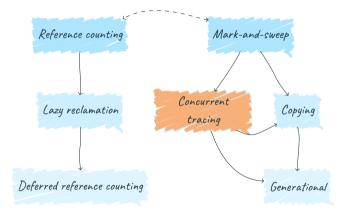




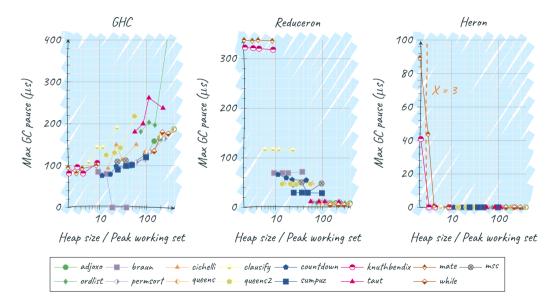




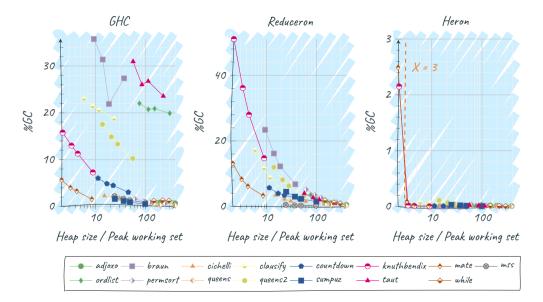




GC worst-case pause



GC wall-clock overhead



Total GC wall-clock pause

